

## PATHS AND CIRCUITS IN $\mathbb{G}$ -GRAPHS OF CERTAIN NON-ABELIAN GROUPS

A. DEWITT\*, A. RODRIGUEZ\*, AND J. DANIEL\*

ABSTRACT. In [BJRTD08], necessary and sufficient conditions were given for the existence of Eulerian and Hamiltonian paths and circuits in the  $\mathbb{G}$ -graph of the dihedral group  $D_n$ . In this paper, we consider the  $\mathbb{G}$ -graphs of the quasihedral, modular, and generalized quaternion group. These groups are of rank 2 and we consider only the graphs  $\Gamma(G, S)$  where  $|S| = 2$ .

### 1. INTRODUCTION

Let  $G$  be a finitely generated group with generating set  $S = \{s_1, \dots, s_k\}$ . For a subgroup  $H$  of  $G$ , define the subset  $T_H$  of  $G$  to be a *left transversal* for  $H$  if  $\{xH \mid x \in T_H\}$  is precisely the set of all left cosets of  $H$  in  $G$ . For each  $s_i \in S$  let  $H_i = \langle s_i \rangle$ . Associate a simple graph  $\Gamma(G, S)$  to  $(G, S)$  with vertex set  $V = \{x_j H_i \mid x_j \in T_{H_i}\}$ . Two distinct vertices  $x_j H_i$  and  $x_l H_k$  in  $V$  are joined by an edge if  $x_j \langle s_i \rangle \cap x_l \langle s_k \rangle$  is nonempty. The edge set  $E$  consists of pairs  $(x_j H_i, x_l H_k)$ .  $\Gamma(G, S)$  defined this way has no multiedge or loop. A multiedge graph was defined similarly in 2004. Many of the results about this graph [[BG04], [BGL05], [BG05], and [BG07]] can be modified for the simple graph,  $\Gamma(G, S)$ , [D08]. The main object of this paper is to study the existence of Eulerian and Hamiltonian paths and circuits in the  $\mathbb{G}$ -graphs of the quasihedral, modular, and generalized quaternion group. To explore the existence of Eulerian paths and circuits in  $\Gamma(G, S)$ , we recall a few theorems of Euler and a result from [BJRTD08].

**Theorem 1.** (Euler) *Let  $\Gamma$  be a nontrivial connected graph. Then  $\Gamma$  has an Eulerian circuit if and only if every vertex is of even degree.*

**Theorem 2.** (Euler) *Let  $\Gamma$  be a nontrivial connected graph. Then  $\Gamma$  has an Eulerian path if and only if  $\Gamma$  has exactly two vertices of odd degree. Furthermore, the path begins at one of the vertices of odd degree and terminates at the other.*

**Lemma 3.** [BJRTD08] *If  $G$  is a group with generating set  $S = \{s_1, \dots, s_n\}$  and  $S_{i,j} = |\langle s_i \rangle \cap \langle s_j \rangle|$ , then the degree of the vertex  $\langle s_i \rangle$ , denoted  $\deg(\langle s_i \rangle)$ , is*

$$\deg(\langle s_i \rangle) = \left( \sum_{j=1}^n |s_i| / S_{i,j} \right) - 1.$$

*Remark 1.* Notice that  $\deg(\langle s_i \rangle) = \deg(x_j \langle s_i \rangle)$  for all  $x_j \langle s_i \rangle$  in  $V_i$ .

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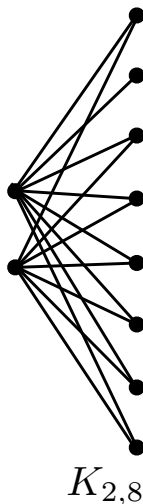
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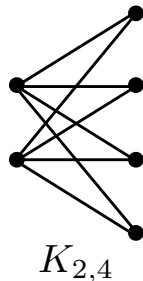
We consider the  $\mathbb{G}$ -graphs of the quasihedral, modular, and generalized quaternion group. We start with a few examples of the graphs.

*Example 1.*

- (i) The modular group,  $M$ , has presentation  $\langle s, t \mid s^8 = t^2 = e, st = ts^5 \rangle$ . Letting  $S = \{s, t\}$ , the  $\mathbb{G}$  graph of this group is  $\Gamma(M, S)$ .



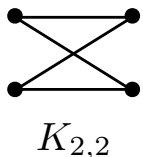
- (ii) The quasihedral group,  $QS$ , has presentation  $\langle s, t \mid s^8 = t^2 = e, st = ts^3 \rangle$ . Letting  $S = \{s, ts\}$ , the  $\mathbb{G}$  graph of this group is  $\Gamma(QS, S)$ .



- (iii) The generalized quaternion group,  $Q_{2^n}$ , has presentation

$$\langle s, t \mid s^{2^{n-1}} = e, s^{2^{n-2}} = t^2, tst^{-1} = s^{-1} \rangle.$$

Letting  $n = 3$ ,  $S = \{s, t\}$ , the  $\mathbb{G}$  graph of this group is  $\Gamma(Q_{2^3}, S)$ .



The next lemma pertains to all of the groups in question.

**Lemma 4.** *Let  $G = M, QS$ , or  $Q_{2^n}$  and let  $j$  be an odd integer then*

$$\langle s^j \rangle = \langle s \rangle = \{s, s^2, \dots, s^{|s|-1}, e\}.$$

*Proof.* For each of the above groups,  $|s|$  is even. So  $\gcd(j, |s|) = 1$  and there exist  $x, y \in \mathbb{Z}$  such that  $jx + |s|y = 1$ . So

$$\begin{aligned} s^1 &= s^{jx+|s|y} \\ s^1 &= s^{jx} s^{|s|y} \\ s^1 &= (s^j)^x (s^{|s|})^y \\ s^1 &= (s^j)^x (e)^y \end{aligned}$$

Therefore  $s^1 = (s^j)^x$  and  $\langle s \rangle = \langle s^j \rangle$ . □

## 2. THE MODULAR GROUP

Recall that the modular group,  $M$ , has presentation  $\langle s, t \mid s^8 = t^2 = e, st = ts^5 \rangle$ . Next we determine the existence of Eulerian and Hamiltonian circuits and paths.

**Lemma 5.** *If  $G$  is the modular group and  $n$  is odd, then*

$$\langle ts^n \rangle = \langle ts \rangle = \{ts, s^6, ts^7, s^4, ts^5, s^2, ts^3, e\}.$$

**Lemma 6.** *If  $G$  is the modular group, then  $\langle ts^2 \rangle = \langle ts^6 \rangle = \{ts^2, s^4, ts^6, e\}$ .*

**Lemma 7.** *If  $G$  is the modular group and  $n = 2$  or  $6$ , then  $|\langle s \rangle \cap \langle ts^n \rangle| = 2$ .*

**Lemma 8.** *If  $G$  is the modular group and  $n$  is odd, then  $|\langle s \rangle \cap \langle ts^n \rangle| = 4$ .*

**Theorem 9.** *If  $G$  is the modular group, and  $S$  is a minimal generating set, then  $\Gamma(G, S)$  contains an Eulerian circuit.*

*Proof.* Let  $G$  be the modular group and  $S$  be a minimal generating set. Then  $S = \{s^n, ts^k\}$ , where  $n$  is odd,  $1 \leq n \leq 7$ , and  $0 \leq k \leq 7$  or  $S = \{ts^n, ts^m\}$ , where  $n$  is odd and  $m$  is even. By using the lemmas above there exists three distinct graphs.

case i) Let  $S = \{s^n, t\}$  where  $n$  is odd, then  $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle t \rangle| = 1$  and  $deg(\langle s^n \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{1} - 1 = 8$ , which is even.

Similarly  $deg(\langle t \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{2}{1} + \frac{2}{2} - 1 = 2$ , which is even. This graph is  $K_{2,8}$  and contains an Eulerian circuit.

case ii) Let  $S = \{s^n, ts^m\}$  where  $n, m$  are odd, then  $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle ts^m \rangle| = 4$  and  $deg(\langle s^n \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{4} - 1 = 2$ , which is

even. Similarly  $deg(\langle ts^m \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{8}{S_{2,1}} + \frac{8}{S_{2,2}} - 1 = \frac{8}{4} + \frac{8}{8} - 1 = 2$ , which is even. This graph is  $K_{2,2}$  and contains an Eulerian circuit.

case iii) Let  $S = \{s^n, ts^k\}$  where  $n$  is odd and  $k = 2$  or  $6$ , then  $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle ts^k \rangle| = 2$  and  $deg(\langle s^n \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{2} - 1 = 4$ ,

which is even. Similarly  $\deg(\langle ts^k \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{4}{2} + \frac{4}{4} - 1 =$

2, which is even. This graph is  $K_{2,4}$  and contains an Eulerian circuit.

case iv) Let  $S = \{s^n, ts^4\}$  where  $n$  is odd, then  $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle ts^4 \rangle| = 1$  and  $\deg(\langle s^n \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{1} - 1 = 8$ , which is even.

Similarly  $\deg(\langle ts^4 \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{2}{1} + \frac{2}{2} - 1 = 2$ , which is even. This graph is  $K_{2,8}$  and contains an Eulerian circuit.

case v) Let  $S = \{ts^n, t\}$  where  $n$  is odd, then  $S_{1,2} = S_{2,1} = |\langle ts^n \rangle \cap \langle t \rangle| = 1$  and  $\deg(\langle ts^n \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{1} - 1 = 8$ , which is even.

Similarly  $\deg(\langle t \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{2}{1} + \frac{2}{2} - 1 = 2$ , which is even. This graph is  $K_{2,8}$  and contains an Eulerian circuit.

case vi) Let  $S = \{ts^n, ts^k\}$  where  $n$  is odd and  $k = 2$  or  $6$ , then  $S_{1,2} = S_{2,1} = |\langle ts^n \rangle \cap \langle ts^k \rangle| = 2$  and  $\deg(\langle ts^n \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{2} - 1 = 4$ , which is even. Similarly  $\deg(\langle ts^k \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{4}{2} + \frac{4}{4} - 1 = 2$ , which is even. This graph is  $K_{2,4}$  and contains an Eulerian circuit.

case vii) Let  $S = \{ts^n, ts^4\}$  where  $n$  is odd, then  $S_{1,2} = S_{2,1} = |\langle ts^n \rangle \cap \langle ts^4 \rangle| = 1$  and  $\deg(\langle ts^n \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{1} - 1 = 8$ , which is

even. Similarly  $\deg(\langle ts^4 \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{2}{1} + \frac{2}{2} - 1 = 2$ , which is even. This graph is  $K_{2,8}$  and contains an Eulerian circuit.  $\square$

*Remark 2.* For all minimal generating sets,  $\Gamma(M, S)$  does not contain an Eulerian path.

**Theorem 10.** *If  $G$  is the modular group, and  $S = \{s^n, ts^m\}$  where  $n, m$  are odd, then  $\Gamma(G, S)$  contains a Hamiltonian circuit and a Hamiltonian path.*

*Proof.* The vertex set of  $\Gamma(M, S)$  is  $V(\Gamma(M, S)) = \{\langle s^n \rangle, t\langle s^n \rangle, \langle ts^m \rangle, t\langle ts^m \rangle\}$ . A Hamiltonian circuit is given by

$$\langle s^n \rangle, \langle ts^m \rangle, t\langle s^n \rangle, t\langle ts^m \rangle, \langle s^n \rangle.$$

A Hamiltonian path is given by

$$\langle s^n \rangle, \langle ts^m \rangle, t\langle s^n \rangle, t\langle ts^m \rangle.$$

$\square$

*Remark 3.*  $S = \{s^n, ts^m\}$  where  $n, m$  are odd is the only minimal generating set of  $M$  that yields a graph that contains a Hamiltonian circuit (path).

### 3. THE QUASIHEDRAL GROUP

Recall that the quasihedral group,  $QS$ , has presentation  $\langle s, t \mid s^8 = t^2 = e, st = ts^3 \rangle$ . Next we determine the existence of Eulerian and Hamiltonian circuits and paths.

**Lemma 11.** *If  $G$  is the quasihedral group and  $n$  is 1 or 5, then  $\langle ts^n \rangle = \{ts, s^4, ts^5, e\}$ .*

**Lemma 12.** *If  $G$  is the quasihedral group and  $n$  is 3 or 7, then  $\langle ts^n \rangle = \{ts^3, s^4, ts^7, e\}$ .*

**Lemma 13.** *If  $G$  is the quasihedral group and  $n$  is even, then  $\langle ts^n \rangle = \{ts^n, e\}$ .*

**Lemma 14.** *If  $G$  is the quasihedral group and  $n$  is even, then  $|\langle s \rangle \cap \langle ts^n \rangle| = 1$ .*

**Lemma 15.** *If  $G$  is the quasihedral group and  $n$  is odd, then  $|\langle s \rangle \cap \langle ts^n \rangle| = 2$ .*

**Theorem 16.** *If  $G$  is the quasihedral group, and  $S$  is a minimal generating set, then  $\Gamma(G, S)$  contains a Eulerian circuit.*

*Proof.* Let  $G$  be the quasihedral group and  $S$  be a minimal generating set. Then  $S = \{s^n, ts^k\}$ , where  $n$  is odd and  $1 \leq n \leq 7$  and  $1 \leq k \leq 3$  or  $S = \{ts^n, ts^m\}$ , where  $n$  is odd and  $m$  is even. By using the above lemmas, there exists three distinct graphs.

case i) Let  $S = \{s^n, ts^m\}$ , where  $n, m$  are odd, then  $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle ts^m \rangle| = 2$  and  $deg(\langle s^n \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{2} - 1 = 4$ , which is

even. Similarly,  $deg(\langle ts^m \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{4}{S_{2,1}} + \frac{4}{S_{2,2}} - 1 = \frac{4}{2} + \frac{4}{4} - 1 = 2$ ,

which is even. This graph is  $K_{2,4}$  and contains an Eulerian circuit.

case ii) Let  $S = \{s^n, ts^m\}$ , where  $n$  is odd and  $m$  is even, then  $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle ts^m \rangle| = 1$  and  $deg(\langle s^n \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{1} - 1 = 8$ ,

which is even. Similarly,  $deg(\langle ts^m \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 =$

$\frac{2}{1} + \frac{2}{2} - 1 = 2$ , which is even. This graph is  $K_{2,8}$  and contains an Eulerian circuit.

case iii) Let  $S = \{ts^n, ts^m\}$  where  $n$  is odd and  $m$  is even, then  $S_{1,2} = S_{2,1} = |\langle ts^n \rangle \cap \langle ts^m \rangle| = 1$  and  $deg(\langle ts^n \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{4}{S_{1,1}} + \frac{4}{S_{1,2}} - 1 = \frac{4}{4} + \frac{4}{1} - 1 =$

$4$ , which is even. Similarly  $deg(\langle ts^m \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 =$

$\frac{2}{1} + \frac{2}{2} - 1 = 2$ , which is even. By applying Euler's theorem, this graph contains an Eulerian circuit.  $\square$

*Remark 4.* For all minimal generating sets,  $\Gamma(QS, S)$  does not contain an Eulerian path, a Hamiltonian path, or a Hamiltonian circuit.

## 4. GENERALIZED QUATERNION GROUP

Recall that the generalized quaternion group,  $Q_{2^n}$ , has presentation  $\langle s, t \mid s^{2^{n-1}} = e, s^{2^{n-2}} = t^2, tst^{-1} = s^{-1} \rangle$ . Next we determine the existence of Eulerian and Hamiltonian circuits and paths.

**Lemma 17.** *If  $G$  is the generalized quaternion group, then  $t^4 = e$ .*

*Proof.* Let  $G$  be the generalized quaternion group. Recall that  $t^2 = s^{2^{n-2}}$ . Squaring both sides,

$$\begin{aligned} (t^2 = s^{2^{n-2}})^2 \\ t^4 = s^{2^{n-1}} = e. \end{aligned}$$

□

**Lemma 18.** *Let  $G$  be the generalized quaternion group, then  $(ts^j)^2 = t^2$  for all  $j$ .*

*Proof.* We proceed with induction on  $j$ . Let  $j = 1$ , then  $(ts^1)^2 = tsts = ts(s^{-1}t) = t^2$  and the theorem holds for  $j = 1$ . Assume that the theorem holds for  $j = k$ , i.e.,  $(ts^k)^2 = t^2$ .

Now let  $j = k + 1$ , then  $(ts^{k+1})^2 = ts^{k+1}ts^{k+1} = ts^{k+1}ts^k = ts^{k+1}s^{-1}ts^k = ts^kts^k = (ts^k)^2 = t^2$ . Therefore  $(ts^j)^2 = t^2$  for all  $j$ .

□

**Lemma 19.** *Let  $G$  be the generalized quaternion group, then  $\langle ts^j \rangle = \{ts^j, t^2, t^3s^j, e\}$  for all  $j$ .*

**Lemma 20.** *If  $G$  is the generalized quaternion group and  $\langle ts^j \rangle \neq \langle ts^k \rangle$ , then  $\langle ts^j \rangle \cap \langle ts^k \rangle = \{t^2, e\}$  and  $|\langle ts^j \rangle \cap \langle ts^k \rangle| = 2$ .*

**Theorem 21.** *If  $G$  is the generalized quaternion group, and  $S$  is a minimal generating set, then  $\Gamma(G, S)$  contains an Eulerian circuit.*

*Proof.* Let  $G$  be the generalized quaternion group and  $S$  be a minimal generating set. Then,  $S = \{s^k, ts^j\}$  where  $k$  is odd or  $S = \{ts^k, ts^m\}$ , where  $k$  is odd and  $m$  is even.

case i) Let  $S = \{s^k, ts^j\}$  where  $k$  is odd, then  $S_{1,2} = S_{2,1} = |\langle s^k \rangle \cap \langle ts^j \rangle| = 2$  and  $deg(\langle s^k \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{2^{n-1}}{S_{1,1}} + \frac{2^{n-1}}{S_{1,2}} - 1 = \frac{2^{n-1}}{2^{n-1}} + \frac{2^{n-1}}{2} - 1 = \frac{2^{n-1}}{2} = 2^{n-2}$ ,

which is even. Similarly  $deg(\langle ts^j \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{4}{S_{2,1}} + \frac{4}{S_{2,2}} - 1 = \frac{4}{4} + \frac{4}{2} - 1 = 2$ ,

which is even. This graph is  $K_{2,2^{n-2}}$  and contains an Eulerian circuit.

case ii) Let  $S = \{ts^k, ts^m\}$ , where  $k$  is odd and  $m$  is even, then  $S_{1,2} = S_{2,1} = |\langle ts^k \rangle \cap \langle ts^m \rangle| = 2$  and  $deg(\langle ts^k \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{4}{S_{1,1}} + \frac{4}{S_{1,2}} - 1 = \frac{4}{4} + \frac{4}{2} - 1 = 2$ ,

which is even. Similarly  $deg(\langle ts^m \rangle) = \left( \sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{4}{S_{2,1}} + \frac{4}{S_{2,2}} - 1 =$

$\frac{4}{4} + \frac{4}{2} - 1 = 2$ , which is even. By applying Euler's theorem, this graph contains an Eulerian circuit.

□

*Remark 5.* For all minimal generating sets,  $\Gamma(Q_{2^n}, S)$  does not contain an Eulerian path.

**Theorem 22.** *If  $G$  is the generalized quaternion group,  $Q_{2^n}$ , and  $S = \{s^k, ts^m\}$  where  $k$  is odd, then  $\Gamma(G, S)$  contains a Hamiltonian circuit and a Hamiltonian path for  $n = 3$ .*

*Proof.* The vertex set of  $\Gamma(Q_{2^3}, S)$  is  $V(\Gamma(Q_{2^3}, S)) = \{\langle s^k \rangle, t\langle s^k \rangle, \langle ts^m \rangle, t\langle ts^m \rangle\}$ . A Hamiltonian circuit is given by

$$\langle s^k \rangle, \langle ts^m \rangle, t\langle s^k \rangle, t\langle ts^m \rangle, \langle s^k \rangle.$$

A Hamiltonian path is given by

$$\langle s^k \rangle, \langle ts^m \rangle, t\langle s^k \rangle, t\langle ts^m \rangle.$$

□

**Theorem 23.** *If  $G$  is the generalized quaternion group,  $Q_{2^n}$ , and  $S = \{ts^k, ts^m\}$ , where  $k$  is odd and  $m$  is even, then  $\Gamma(G, S)$  contains a Hamiltonian circuit and a Hamiltonian path.*

*Proof.* The vertex set of  $\Gamma(Q_{2^n}, S)$  is

$$V(\Gamma(Q_{2^n}, S)) = \{\langle ts^k \rangle, s\langle ts^k \rangle, \dots, s^{2^{n-2}-1}\langle ts^k \rangle, \langle ts^m \rangle, s\langle ts^m \rangle, \dots, s^{2^{n-2}-1}\langle ts^m \rangle\}.$$

A Hamiltonian circuit is given by

$$\langle ts^k \rangle, \langle ts^m \rangle, s^{k-m}\langle ts^k \rangle, s^{k-m}\langle ts^m \rangle, \dots, s^{k-(2^{n-2}-1)m}\langle ts^k \rangle, s^{k-(2^{n-2}-1)m}\langle ts^m \rangle, \langle ts^k \rangle.$$

A Hamiltonian path is given by

$$\langle ts^k \rangle, \langle ts^m \rangle, s^{k-m}\langle ts^k \rangle, s^{k-m}\langle ts^m \rangle, s^{k-2m}\langle ts^k \rangle, s^{k-2m}\langle ts^m \rangle, \dots, s^{k-(2^{n-2}-1)m}\langle ts^k \rangle, s^{k-(2^{n-2}-1)m}\langle ts^m \rangle.$$

□

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DEPARTMENT OF MATHEMATICS, LAMAR UNIVERSITY, BEAUMONT, TX 77710  
*E-mail address:* aldewitt@my.lamar.edu

DEPARTMENT OF MATHEMATICS, LAMAR UNIVERSITY, BEAUMONT, TX 77710  
*E-mail address:* amrodriguez1@my.lamar.edu

DEPARTMENT OF MATHEMATICS, LAMAR UNIVERSITY, BEAUMONT, TX 77710  
*E-mail address:* Jennifer.Daniel@lamar.edu