CHAPTER 1: FUN AND GAMES

Finite Mathematics Spring 2017

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- The winner is the last player to remove a token.



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- $\textcircled{\textbf{0}} \hspace{0.1in} \textbf{S} = \{\textbf{1}, \textbf{2}, \textbf{3}\} \hspace{0.1in} \text{and} \hspace{0.1in}$
 - **n** = 4
 - **n** = **9**
 - Do you want to be Player 1 or Player 2?

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Can you think of a way to determine the winner and a strategy for that player for any given **S** and **n**?

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- The key idea is to make up these sets recursively.

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Example: $\mathbf{S} = \{\mathbf{1}, \mathbf{2}, \mathbf{5}\}$

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- Put **3** into \mathcal{P} . Why?

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- Put 3 into \mathcal{P} . Why?
- Put 4 into $\mathcal N$, and then put 6 into $\mathcal P$. Why?

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Practice this strategy

Example: $S = \{1, 3, 4\}$ Question: Can you do $S = \{1, 2, 6, 9\}$?

Summary

- **(**) We build the sets \mathcal{N} and \mathcal{P} recursively.
- **2** Put 0 and each integer from **S** in \mathcal{N} .
- Consider the smallest positive integer k that has not been put into either N or P.
- If there is at least one number in S so that when that many tokens are removed from a pile of k tokens the number of tokens remaining in the pile is in P, then put k into N.
- If the above is not true, then for every number in S when that many tokens are removed from a pile of k tokens the number of tokens remaining in the pile is in N. Put k into P.
- O Repeat starting at step 3 above.