Can a Machine Think? Mathematics 15: Lecture 27

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Can a Machine Think?

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- ► Turing committed suicide in 1954.

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Turing shows that such a machine could perform any possible computation.

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- But T would not halt for any other input.

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- But now A(n, n) is a Turing machine: it's a computational procedure which takes an input n and may or may not stop with some output. So A(n, n) = T_k(n) for some k.
- We now have: If A(k, k) stops, then T_k(k) does not stop; that is, if T_k(k) stops, then T_k(k) does not stop.

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- This means that we know T_k(k) does not stop, but no computational procedure can tell us that it does.

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- Machine: programmable digital computer, in part because it can mimic any other type of machine.
- Thinking: we must say a machine can think if an interrogator cannot distinguish the replies from the machine and a person.

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- Extra-sensory perception